

From Program Execution to Automatic Reasoning

Integrating Ontologies into Programming Languages

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SYMPOSIUM ON LANGUAGES,
APPLICATIONS AND TECHNOLOGIES

State of the Art

- Widely used object-oriented programming languages include a built-in static type system
- Conceptual framework that makes it particularly easy to
 - design,
 - understand, and
 - maintain object-oriented systems
- In recent years *schema languages* and *ontology languages* emerged
 - Programming language agnostic data types and content models
 - Automatic reasoning about domain model
- Powerful modeling features readily available
- Impedance mismatch makes programming difficult and error-prone
- Facilitate the development of intelligent applications

The Extensible Markup Language (XML)

- Separate formatting of a document from its content
- *General Markup Language* (GML) proposed by Charles Goldfarb, Edward Mosher, and Raymond Lorie in 1967
- Standard General Markup Language (SGML) became an ISO standard in 1986
- In 1999, the Extensible Markup Language (XML) became a W3C standard
- XML is not a language

XML Schema Definition (XSD)

- Published as a W3C recommendation in 2001
- XSD specification comprises two parts
 - XML Schema Part 1: Structures
 - XML Schema Part 2: Datatypes
- Atomic XSD data types are triples
 - Value space
 - Lexical space
 - Fundamental facets

Built-in XSD Data Types

- 19 primitive built-in data types
- Several isomorphic mappings to programming language data types

.NET value type	XSD type	.NET value type	XSD type
<i>System.String</i>	<i>xsd:string</i>	<i>System.Int16</i>	<i>xsd:short</i>
<i>System.Boolean</i>	<i>xsd:boolean</i>	<i>System.UInt16</i>	<i>xsd:unsignedShort</i>
<i>System.Single</i>	<i>xsd:float</i>	<i>System.Int32</i>	<i>xsd:int</i>
<i>System.Double</i>	<i>xsd:double</i>	<i>System.UInt32</i>	<i>xsd:unsignedInt</i>
<i>System.SByte</i>	<i>xsd:byte</i>	<i>System.Int64</i>	<i>xsd:long</i>
<i>System.Byte</i>	<i>xsd:unsignedByte</i>	<i>System.UInt64</i>	<i>xsd:unsignedLong</i>

- Why are there no isomorphic mappings in general?

User-defined Type „age“

- A valid age is any integer number greater than or equal to zero and less than 110.

```
<xsd:schema xmlns:xsd=http://www.w3.org/2001/XMLSchema>  
  <xsd:simpleType name="age">  
    <xsd:restriction base="xsd:int">  
      <xsd:minInclusive value="0" />  
      <xsd:maxExclusive value="110" />  
    </xsd:restriction>  
  </xsd:simpleType>  
</xsd:schema>
```

XSD Data Types in C#

- Mapping option 1

```
class Person {  
    public uint HasAge;  
}
```

- Mapping option 2

```
using System.Diagnostics;  
internal class Person {  
    private uint _HasAge;  
    public uint HasAge {  
        get { return _HasAge; }  
        set { Trace.Assert(value >= 0 && value < 110);  
            _HasAge = value; }  
    }  
}
```

XSD Data Types in C# cont'd

- Mapping option 3

```
using System.Diagnostics;
internal struct uint110 {
    private uint value;
    public static implicit operator uint110(uint value) {
        Trace.Assert(value < 100);
        return new uint110 {value = value};
    }
    public static implicit operator uint(uint110 value) {
        return value.value;
    }
}
```


XSD Data Types in C# cont'd

- Mapping option 4

```
abstract class XsdAnySimpleType {  
    protected object value;  
}  
class XsdUnsignedInt : XsdAnySimpleType {  
    public static implicit operator  
        XsdUnsignedInt(uint value) {  
        return new XsdUnsignedInt {value = value};  
    }  
    public static implicit operator  
        uint (XsdUnsignedInt value) {  
        return (uint) value.value;  
    }  
}
```

XSD Constraining Facets

- 12 constraining facets

Constraint	Description
<i>length</i>	defines the number of <i>units of length</i>
<i>minLength</i>	defines the minimum number of <i>units of length</i>
<i>maxLength</i>	defines the maximum number of <i>units of length</i>
<i>pattern</i>	constrains the lexical space to literals that match a specific pattern
<i>enumeration</i>	constrains the value space to a specified set of values
<i>minInclusive</i>	defines the inclusive lower bound of the value space
<i>minExclusive</i>	defines the exclusive lower bound of the value space
<i>maxExclusive</i>	defines the exclusive upper bound of the value space
<i>maxInclusive</i>	defines the inclusive upper bound of the value space
<i>totalDigits</i>	defines the maximum number of values in the value space
<i>fractionDigits</i>	defines the minimum difference between values in the value space
<i>whiteSpace</i>	controls the normalization of string data types (modifying)

- A *constraining facet* is an optional property that can be applied to an atomic data type to constrain its value space
- A value space $v(T)$ is the set of values for a given atomic data type T

Constraint Application

- The value space function $v(P)$ is defined semantically for each primitive base type P under consideration
- The value space of a base type T can be restricted by the application of one or more constraints $c_i = \phi(TV)b_i$ $i \in 1..n$
- A constraint body $b = \{x | x \in v(TV)\} \cap \bigcap_{k \in 1..m} \{x | x \prec literal_k\}$ defines the intersection of the value space of TV and those values that satisfy the properties $x \prec literal_k$ $k \in 1..m$
- $T_n = T.c_1 \dots c_n \equiv \bigcap_{i \in 1..n}^T c_i$

Constraint Subsumption

- S-CSTRVSPACE (non-algorithmic)

$$\frac{v(c_1\{\{TV \leftarrow T\}\}) \subseteq v(c_2\{\{TV \leftarrow T\}\})}{c_1 <:: c_2}$$

- S-CSTRWIDTH

$$\phi(TV)\{x|x \in v(TV)\} \bigcap_{i \in 1..n+k} \{x|x \prec y_i\} <:: \phi(TV)\{x|x \in v(TV)\} \bigcap_{i \in 1..n} \{x|x \prec y_i\}$$

- S-CSTRDEPTH

$$\frac{\text{for each } i \{x|x \prec y_i\} \subseteq \{x|x \prec z_i\}}{\phi(TV)\{x|x \in v(TV)\} \bigcap_{i \in 1..n} \{x|x \prec y_i\} <:: \phi(TV)\{x|x \in v(TV)\} \bigcap_{i \in 1..n} \{x|x \prec z_i\}}$$

Subtyping of XSD Data Types

- S-VSPACE (non-algorithmic)

$$\frac{v(S) \subseteq v(T)}{S <: T}$$

- S-WIDTH

$$\frac{S = \bigcap_{i \in 1..n+k}^T c_i \quad U = \bigcap_{i \in 1..n}^T c_i}{S <: U}$$

- S-DEPTH

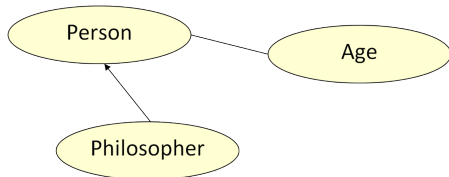
$$\frac{\text{for each } i \quad c_i <:: d_i \quad S = \bigcap_{i \in 1..n}^T c_i \quad U = \bigcap_{i \in 1..n}^T d_i}{S <: U}$$

An XSD Type System

- Interpretation of XSD data types as computational structures
- Foundation for the integration into programming language type system
- Will it be *sound* (will there be *progress* and *preservation*)?
- Yes, it will because...
 - XSD type construction does not pertain to programming language typing and evaluation rules.
 - Subsumption property (if $S' <: S$ and $S <: T$ then $S' <: T$) can be easily proved.
- Practical implementation in the Zhi# programming language

Initial Knowledge Representation Schemes

- In the 1960s, network-based notations occurred
- Hierarchical structures in monotonic inheritance networks (1970s)



- Translation of semantic networks into fragments of first-order predicate calculus (1980s)

Philosophers are men.

Therefore philosophers' thoughts are thoughts of men.

Description Logics (DL)

- Decision procedures of DL systems always terminate
- Efficient query answering
- Define the concepts of a domain of discourse
- Use concepts to specify properties of objects
- General knowledge contained in the *TBox*
- Contingent knowledge contained in the *ABox*
- Fundamental inference is *subsumption*: $C \sqsubseteq D$

The \mathcal{AL} Description Logic

- Basic Description Logic introduced 1991

Constructor Name	Syntax	Semantics
atomic concept	A	$A^{\mathcal{I}} \subseteq \Delta^{\mathcal{I}}$
top level concept	\top	$\top^{\mathcal{I}} = \Delta^{\mathcal{I}}$
bottom concept	\perp	$\perp^{\mathcal{I}} = \{\}$
atomic negation	$\neg A$	$(\neg A)^{\mathcal{I}} = \Delta^{\mathcal{I}} \setminus A^{\mathcal{I}}$
intersection	$C_1 \sqcap C_2$	$(C_1 \sqcap C_2)^{\mathcal{I}} = C_1^{\mathcal{I}} \cap C_2^{\mathcal{I}}$
value restriction	$\forall R.C$	$(\forall R.C)^{\mathcal{I}} = \{x \mid \forall y. \langle x, y \rangle \in R^{\mathcal{I}} \rightarrow y \in C^{\mathcal{I}}\}$
ltd. ex. quantification	$\exists R.\top$	$(\exists R.\top)^{\mathcal{I}} = \{x \mid \exists y. \langle x, y \rangle \in R^{\mathcal{I}}\}$

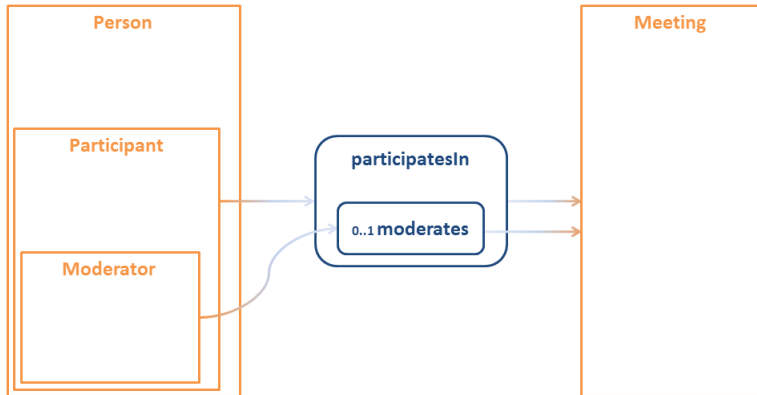
- Set theoretic interpretation $\mathcal{I} = (\Delta^{\mathcal{I}}, \cdot^{\mathcal{I}})$

A Business Meeting Scenario

- Atomic concepts: *Person*, *Meeting*
- General concept inclusion: *Employee* \sqsubseteq *Person*
- Ltd. ex. quantification:
 $MeetingParticipant \equiv Person \sqcap \exists attendsMeeting.\top$
- Value restriction:
 $MarketingMeeting \equiv Meeting \sqcap \forall hasTopic.MarkingTopic$
- Negation: $Guest \equiv Person \sqcap \neg Employee$

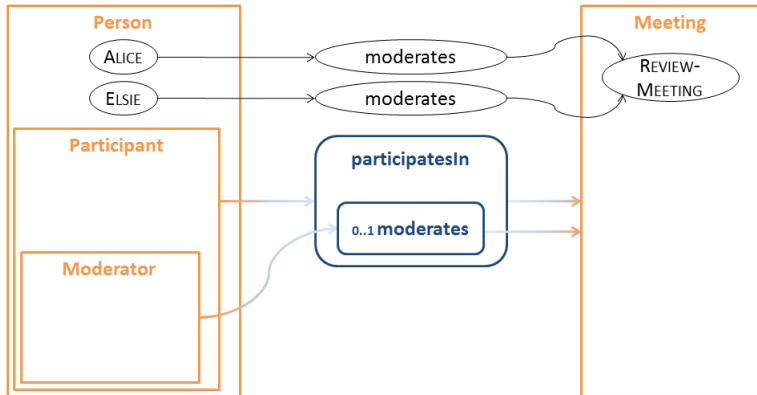
Reasoning With The *SHOIN* DL

- *SHOIN* extends \mathcal{AL} with additional modelling features: negation (\mathcal{C}), qualified number restrictions (\mathcal{Q}), role hierarchies (\mathcal{H}), and inverse (\mathcal{I}) and transitive roles (\mathcal{R}_+)
- *SHOIN(D)* includes a concrete domain: data types



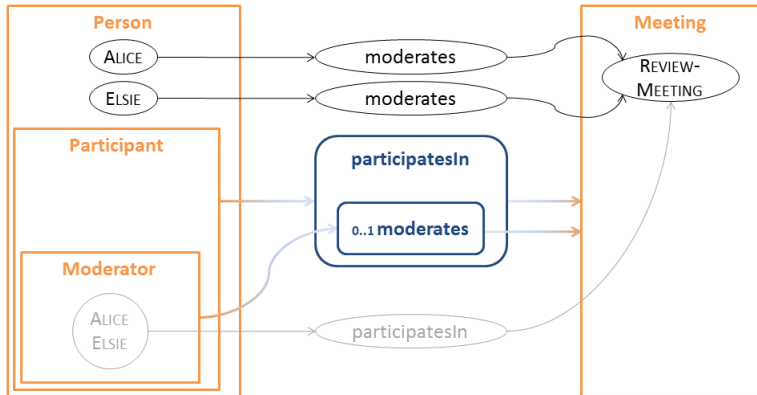
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API-based Ontology Management

- No programming language integration
- Insufficient compiler/IDE support
- Jena Semantic Web Framework widely used

```

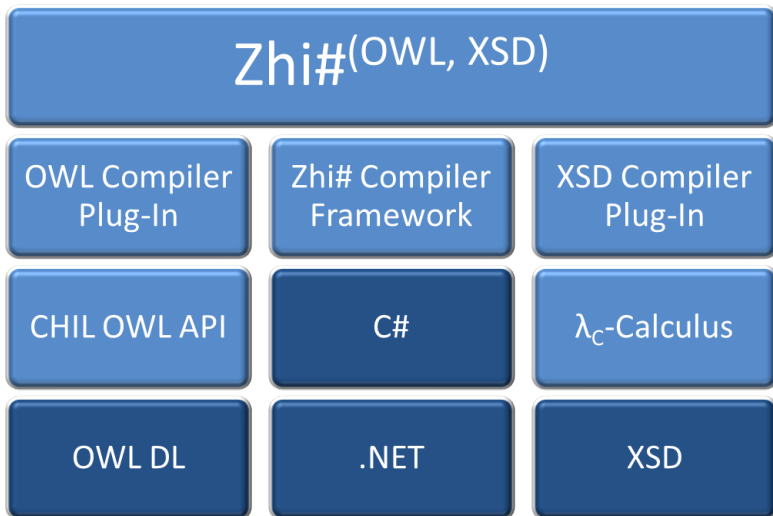
1 public class Program {
2     public static void main(String [] args) {
3         String eval = "http://www.zhimantic.com/eval#";
4
5         OntModel m = ModelFactory.createOntologyModel(PelletReasonerFactory.THE_SPEC);
6         m.read(new FileInputStream("Evaluation.owl"), "");
7
8         Individual a = m.getOntClass("http://www.w3.org/2002/07/owl#Thing").
9             createIndividual(eval + "a");
10
11        Individual o = m.getOntClass("http://www.w3.org/2002/07/owl#Thing").
12            createIndividual(eval + "o");
13
14        a.addProperty(m.getObjectProperty(eval + "R"), o);
15
16        for (Iterator it = m.getOntClass(eval + "B").listInstances(); it.hasNext();) {
17            System.out.println(((Individual) it.next()).getURI());
18        }
19
20        o.addProperty(m.getDatatypeProperty(eval + "T"),
21            m.createTypedLiteral("23", "http://www.w3.org/2001/XMLSchema#positiveInteger"));
22    }
23 }

```

Ontologies vs. Object-Orientation

- In Java/C# properties are class members
- In DL ontologies. . .
 - . . . properties form a separate hierarchy (i.e. property centric modeling)
 - . . . property domain and range restrictions are used for ontological reasoning
- Description Logics make the *open world assumption*
- Description Logics do not make the *unique name assumption*

The Zhi# („Semantic C#“) Approach



XSD Data Types in Zhi# Programs

- Precise static typing and type inference for XSD data types

```

1  using System;
2  import XML xsd = http://www.w3.org/2001/XMLSchema;
3  import XML app = http://www.example.com/datatypes;
4  class C {
5      public void f() {
6          int i = new Random().Next(); //i:int
7          if (i < 110)
8              #app#age a = i; //a:int{>= 0}{< 110}, i:int{< 110} Error!
9          }
10 }

```

- Implement OCL expressions with XSD-like data types

Person
-name : String
-age : Integer
-eMail : String

context Person

inv age >= 0 && age < 110

context Person

inv name.size() < 40

Ontologies in Zhi# Programs

- Combination of static typing and dynamic checking
- Full support for datatype properties

```

1 using System;
2 import XML xsd = http://www.w3.org/2001/XMLSchema;
3 import OWL ont = http://www.zhimantic.com/ont;
4 class C {
5     public DateTime f(#ont#Meeting m, #xsd#dateTime xdt) {
6         m.#ont#beginsAt = xdt;
7         #xsd#duration xd = "P0Y0M0DT1H30M0S";
8         m.#ont#endsAt = xdt + xd;
9
10        #ont#Person alice = new #ont#Person(" Alice");
11        #ont#Person bob = new #ont#Person(" Bob");
12        #ont#Person charlie = new #ont#Person(" Charlie");
13
14        m.#ont#hasParticipant = bob;
15        charlie.#ont#participatesIn = m;
16        alice.#ont#moderates = m;
17
18        foreach (#ont#Participant p in m.#ont#hasParticipant)
19            Console.WriteLine(p + " participates in " + m);
20
21        if (m is #ont#ImportantMeeting)
22            Console.WriteLine(m + " is an important meeting");
23
24        return (DateTime) m.#ont#endsAt;
25    }
26 }

```

Conclusion

- Schema and ontology languages can be integrated in object-oriented programming language
- Object-oriented notation sufficient
- Separation of concerns without impedance mismatch
- Applications become intelligent
- Pay as you go
- Programming languages need to become aware of information spaces
- First Workshop on Programming the Semantic Web
www.inf.puc-rio.br/~psw12/

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The Zhi# Approach



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